Perfect simulation without looking back. A case study using a Gibbs sampler.

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Introduction - Markov chains

• A Markov chain, written $X_1, X_2, X_3, \dots \in E$ is the stochastic analogue of a discrete time dynamical system:

$$X_{t+1} = F(X_t).$$

- If F(x) is fixed (deterministic) we have a dynamical system
- If F(x) is chosen randomly at each iteration, but with some fixed transition density

$$\mathbb{P}(F(x) = y) = p(x, y),$$

then we have a Markov chain.

• Often only p(x, y) is specified - then there are many possible choices of F (Borovkov + Foss, 1992)

Introduction - ergodicity

- A Markov chain which visits all parts of its state space *E* sufficiently frequently is called positive recurrent
- In this case, if it is not periodic, the chain settle down over time to some equilibrium distribution π on E:

$$\lim_{t \to \infty} \mathbb{P}(X_t \in A) = \pi(A)$$
$$\lim_{t \to \infty} \frac{1}{t} \sum_{s=1}^t f(X_s) = \int f d\pi \text{ with probability 1.}$$

 Both these results are interesting for Statistics (and Physics, and ...) as a way of doing approximate integration over a given probability distribution π in a complicated space.





Read Once CFTP: a theorem

- $P(x, dy) = \mathbb{P}(F(x) \in dy).$
- Assume $\mathbb{P}(F \text{ is coalescent}) = \epsilon > 0$, and write $\mu(dy) = \mathbb{P}(F(x_0) \in dy \mid F \text{ is coalescent})$
- This gives $P(x, dy) = (1 \epsilon)Q(x, dy) + \epsilon \mu(dy)$, and $\pi P = \pi$.
- **Theorem**. We have

$$\pi = \epsilon \sum_{s=0}^{\infty} (1-\epsilon)^s \mu Q^s.$$

Proof of Theorem

• Theorem. We have

$$\pi = \epsilon \sum_{s=0}^{\infty} (1-\epsilon)^s \mu Q^s.$$
 (1)

• **Proof** Using stationarity $\pi P = \pi$, we have

$$(1-\epsilon)^{k}\pi Q^{k} = (1-\epsilon)^{k-1}\pi (P-\epsilon\mu)Q^{k-1}$$

= $(1-\epsilon)^{k-1}\pi Q^{k-1} - \epsilon(1-\epsilon)^{k-1}\mu Q^{k-1}$
= \cdots =
$$= \pi - \epsilon \sum_{s=1}^{k} (1-\epsilon)^{k-s}\mu Q^{k-s}.$$

As identities between positive kernels, these are true when applied to any bounded test function $f: E \to \mathbb{R}$. Changing variables $k - s \to s$ gives

$$(1-\epsilon)^k \langle \pi Q^k, f \rangle = \langle \pi, f \rangle - \epsilon \sum_{s=0}^{k-1} (1-\epsilon)^s \langle \mu Q^s, f \rangle.$$

Since also $|\langle \pi Q^k, f \rangle| \leq ||f||$, it is now clear that we obtain (1) by letting $k \to \infty$.

Ways of coupling Markov chains I

- Very easy when state space is finite: $E = \{0, ..., n\}$.
- •
- $\begin{cases} \mathbb{P}(X_{t+1} = X_t \pm 1) = 1/2 & \text{if } 0 < X_t < n \\ \mathbb{P}(X_{t+1} = n 1) = 1 & \text{if } X_t = n \\ \mathbb{P}(X_{t+1} = 1) = 1 & \text{if } X_t = 0. \end{cases}$
- Here π is shte uniform distribution on $\{0, \ldots, n\}$.
- Take X_t, X'_t to be independent copies of this chain.
- Can show that $\mathbb{P}(T < \infty) = 1$, so we have convergence!

Ways of coupling Markov chains II

- For general state spaces E, use the Splitting technique. (Nummelin, Athreya + Ney, Meyn + Tweedie, Rosenthal)
- Let X_t and X'_t be two independent chains with transition density p(x, y). Assume there exists a set $C \subset E$ (called a small set!) such that

$$\min_{x \in C} p(x, y) \ge \epsilon \mu(y), \qquad \int \mu(y) dy = 1.$$

- Wait until both chains are simultaneously in C at some time τ say.
- With probability ϵ , choose $Y \sim \mu$ and make both chains jump to Y, i.e. $X_{\tau+1} = X'_{\tau+1} = Y$.
- With probability 1ϵ , set

$$X_{\tau+1} \sim Q(X_{\tau}, \cdot) = (1-\epsilon)^{-1} (p(X_{\tau}, \cdot) - \epsilon \mu(\cdot))$$
$$X'_{\tau+1} \sim Q(X'_{\tau}, \cdot) = (1-\epsilon)^{-1} (p(X'_{\tau}, \cdot) - \epsilon \mu(\cdot))$$

Coupling without Analysis

- Normally at each iteration, we use a randomly generated function $F_t(x)$ such that $X_{t+1} = F_t(X_t)$.
- If $F_t(x) = F_t(x')$ holds for some x, x', then there is the *possibility* of coupling.
- Strategy: From F, generate a new random function $C_Y(F)$ which has a higher chance of coupling.
- **Definition:** Let Y be independent of F, with $\mathbb{P}(Y = y) = q(y)$ say.

$$\mathcal{C}_{Y}(F)(x) = \begin{cases} Y & \text{if } \frac{p(x,Y)q(F(x))}{p(x,F(x))q(Y)} > U[0,1] \\ F(x) & \text{otherwise.} \end{cases}$$

Theorem I

• Suppose the state space E can be entirely covered by a finite union of 1-small sets C_1, C_2, \ldots, C_N say, and suppose we can choose a probability density q in such a way that

 $\epsilon_i \mu_i(\cdot) \le \inf_{z \in C_i} p(z, \cdot) \le \sup_{z \in C_i} p(z, \cdot) \le \gamma_i q(\cdot), \quad i = 1, \dots, N,$

where the γ_i are constants. Then if we make the proposals Y_1, Y_2, \ldots , we have

$$\mathbb{P}\Big(\mathcal{C}_{Y_t} \circ \cdots \circ \mathcal{C}_{Y_1}(F) \text{ has finite range eventually}\Big) = 1.$$

• The convergence is geometrically fast!

Theorem II

- We are never sure that we have made enough proposals Y_1, Y_2, \ldots to guarantee that the new map $\mathcal{C}_{Y_n} \circ \cdots \circ \mathcal{C}_{Y_1}(F)$ has finite range. Fortunately, there exists a different procedure which allows us to generate a finitely coupled map with certainty.
- (Theorem) By using the "Dead Leaves Method" we can generate in a finite number of iterations a finitely coupled map with certainty!

Theorem III

 Suppose we can couple the transition maps F_t(x) into new maps F̃_t(x) which have finite range with probability 1. This only works if the transition density p(x, y) is uniformly ergodic. Then any two Markov chains X̃_t and X̃'_t defined by

 $\widetilde{X}_{t+1} = \widetilde{F}_t(\widetilde{X}_t), \quad \widetilde{X}'_{t+1} = \widetilde{F}_t(\widetilde{X}'_t), \quad \widetilde{X}_0, \widetilde{X}'_0 \text{ arbitrary}$

will couple successfully in a finite time. Moreover, *all* Markov chains of the above type, where \tilde{X}_0 ranges over all points in the state space, must coalesce in a finite time!

Proof of Theorem I

Consider the basins of attraction

$$Basin(Y, F, \xi) = \Big\{ x : p(x, Y)q(F(x)) > \xi p(x, F(x))q(Y) \Big\}.$$









 $C_{Y_2}(C_{Y_l^{(f)}})$



 $C_{Y_5}(C_{Y_4}(C_Y(C_Y(C_Y(f)))))$

Proof of Theorem II

Consider the core regions in each basin

$$\operatorname{Core}(Y,\xi) = \left\{ x : x \in \operatorname{Basin}(Y,F,\xi) \text{ for all } F \right\}$$







Assumptions for Read-Once CFTP

Assumption I A source of independent random update functions f(x) exists, each satisfying

$$\int \pi(x) \mathbb{P}\Big(f(x) \in dy\Big) dx = \pi(y) dy,$$

and the simulation of π should use updates of this type. The target density $\pi(x)$ is known up to a normalization constant.

Assumption II The update functions f have a common probability density p(x, y), which is known up to a normalizing constant:

$$\mathbb{P}\Big(f(x) \in dy\Big) = p(x, y)dy, \quad x \in E.$$

Building the maps F_t

(a) Resetting the state Let b(x) be a proposal density, $\pi(x)/b(x) \to 0$ as $x \to \infty$

$$R_B(x) = \begin{cases} B & \text{if } \pi(B)b(x) > \psi\pi(x)b(B) \\ x & \text{otherwise,} \end{cases}$$

where $\psi \sim U[0,1]$.

(b) Coupling updates f_1, f_2, \ldots Given $x \mapsto f(x)$ (Assumption I), write

$$\mathcal{C}_{Y}(f)(x) = \begin{cases} Y & \text{if } p(x,Y)q(f(x)) > \xi p(x,f(x))q(Y) \\ f(x) & \text{otherwise,} \end{cases}$$

where $\xi \sim U[0, 1]$ independently, and p(x, y) comes from Assumption II. Choose a finite IID sequence $\mathcal{Y} = (\mathcal{Y}_{\infty}, \mathcal{Y}_{\in}, \dots, \mathcal{Y}_{\tau})$ from q. $\mathcal{C}_{\mathcal{Y}}(f) := \mathcal{C}_{Y_1, Y_2, \dots, Y_{\tau}}(f) = \mathcal{C}_{Y_{\tau}} \circ \mathcal{C}_{Y_{\tau-1}} \circ \dots \circ \mathcal{C}_{Y_2} \circ \mathcal{C}_{Y_1}(f).$

(c) Definition of maps F_t Choose f_1, \ldots, f_m , and put

$$F(x) = \mathcal{C}_{\mathcal{Y}_m}(f_m) \circ \cdots \circ \mathcal{C}_{\mathcal{Y}_1}(f_1) \circ R_B(x),$$

Effect of Independence Sampler

Which is more likely? For each pair (x, B), the map

$$R_B(x) = \begin{cases} B & \text{if } \pi(B)b(x) > \psi\pi(x)b(B) \\ x & \text{otherwise,} \end{cases}$$

chooses the most likely configuration



Pump Example (Autogamma)

• Let $x = (\beta, \lambda_1, \dots, \lambda_{10})$, simulate

$$\pi(x) = \exp\left\{ (10\alpha + \gamma - 1) \log \beta - \delta\beta + \sum_{k=1}^{10} \left((s_k + \alpha - 1) \log \lambda_k - (\beta + t_k)\lambda_k \right) \right\},\$$

- Gibbs sampler: One sweep is
 - $f: (\beta, \lambda_1, \dots, \lambda_{10}) \mapsto (\beta', \lambda'_1, \dots, \lambda'_{10})$, where

$$\beta' \sim \pi_0(\cdot \mid \lambda_1, \dots, \lambda_{10}) = \Gamma(\gamma + 10\alpha, \delta + \sum_{k=1}^{10} \lambda_k),$$
$$\lambda'_k \sim \pi_k(\cdot \mid \beta') = \Gamma(\alpha + s_k, \beta' + t_k), \quad k = 1, \dots, 10.$$

• Transition density is

$$p(\beta, \lambda_1, \dots, \lambda_{10}; b, l_1, \dots, l_{10}) = \pi_0(b \,|\, \lambda_1, \dots, \lambda_{10}) \prod_{k=1}^{10} \pi_k(l_k \,|\, b).$$

Pump Example: Constructing $R_B(x)$

• Take $B = (B_0, \dots, B_{10})$ such that $B_0 \sim \Gamma(\gamma, \delta)$, $B_k \sim \Gamma(\alpha, B_0)$ for $k \ge 1$. Then

$$b(x)/\pi(x) = \Gamma(\gamma)^{-1}\Gamma(\alpha)^{-10}\delta^{\gamma} \exp\left(-\sum_{k=1}^{10} (s_k \log \lambda_k - t_k \lambda_k)\right).$$

• Consequently,

$$K_B = \Big\{ x : \sum_{k=1}^{10} (s_k \log \lambda_k - t_k \lambda_k) \ge -\log \psi \\ + \sum_{k=1}^{10} (s_k \log B_k - t_k B_k) \Big\}.$$

• Simpler to take $|\lambda| = \lambda_1 + \cdots + \lambda_{10}$,

$$K_B \subset \left\{ (\beta, \lambda_1, \dots, \lambda_{10}) : 0 \le |\lambda| \le \left(\log \psi - \sum_{k=1}^{10} (s_k \log B_k - t_k B_k) \right) / \max_j t_j \right\}.$$

Pump Example: Constructing $C_Y(f)$

• Take $Y \sim q$ where

$$q(y_0, y_1, \dots, y_{10}) = \pi_0(y_0 \,|\, \lambda_1^*, \dots, \lambda_{10}^*) \prod_{k=1}^{10} \pi_k(y_k \,|\, y_0),$$

• After simplification,

$$\mathcal{C}_{Y}(f)(x) = \begin{cases} Y & \text{if } \exp\left((\beta' - Y_{0})(|\lambda| - |\lambda^{*}|)\right) > \xi, \\ (\beta', \lambda'_{1}, \dots, \lambda'_{10}) & \text{otherwise.} \end{cases}$$

• Thus

$$\operatorname{Basin}(Y, f, \xi) = \Big\{ x : \Big(\psi_0 - (\delta + |\lambda|) Y_0 \Big) \Big(|\lambda| - |\lambda^*| \Big) > \log \xi \Big\},\$$

• After simplification,

$$\operatorname{Basin}(Y, f, \xi) = \left\{ x : |\lambda| \in \left[\frac{-b - \sqrt{b^2 - 4ac}}{2a}, \frac{-b + \sqrt{b^2 - 4ac}}{2a} \right] \right\},$$

where $a = Y_0, b = \log \xi - Y_0(|\lambda^*| - \delta)$ and $c = \delta(\log \xi - Y_0 |\lambda^*|) - \psi_0.$